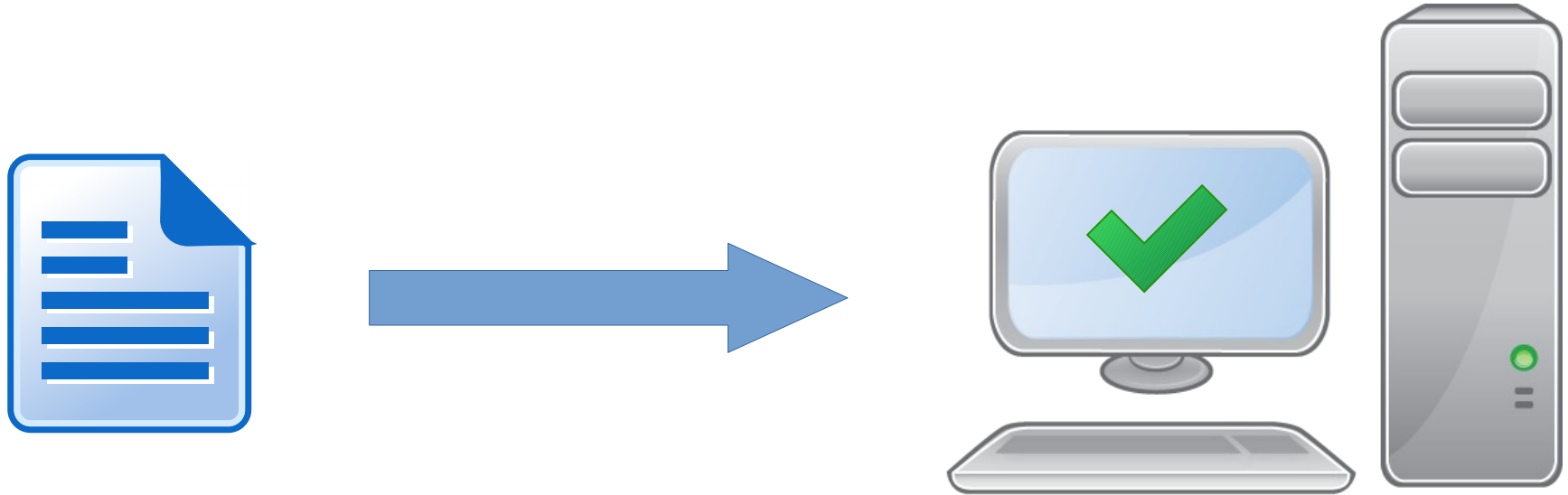


# Verifying C Data-types for Cogent

Louis Francis Cheung (z5062193)

# How Programs Should Work



# What Normally Happens



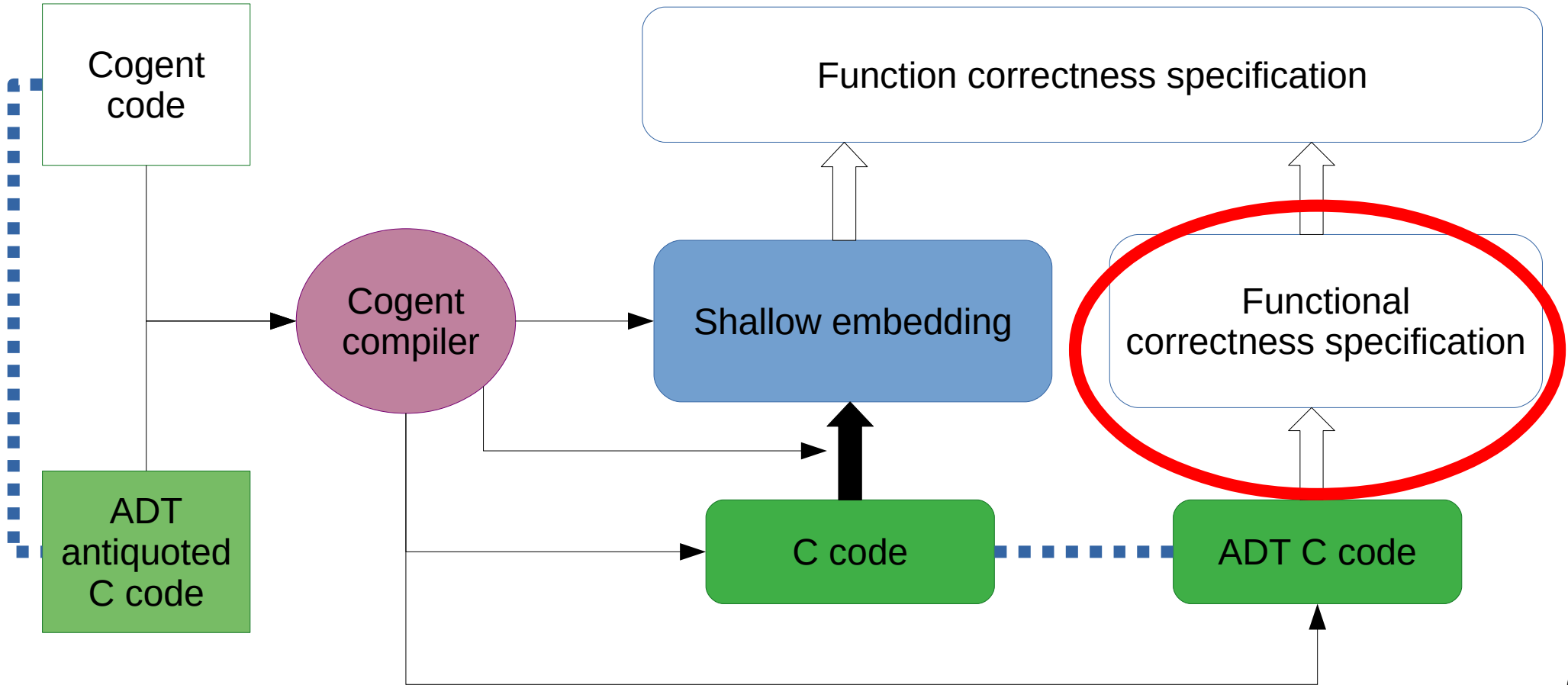
# Verification = Trust



# Cogent

- Restricted, pure, polymorphic, functional language with linear types and no runtime environment nor garbage collector
- Has a certifying compiler
  - Isabelle/HOL shallow embedding
  - C code
  - Refinement proof
- No native support for recursion
- Use FFI to call C functions that implement ADTs and iterators
  - The C functions need to be verified manually

# The Big Picture



# Expected Outcomes I

- Define a functional correctness specification for 32-bit word arrays
- Verify functional correctness of the 32-bit word array implementation
- Verify that the 32-bit word array implementation satisfies Cogent's frame constraints

# Expected Outcomes II

- Formulate a generic specification for word arrays
- Formulate generic proof techniques and requirements to prove functional correctness of word arrays
- Formulate generic proof techniques and requirements to prove frame constraint satisfiability of word arrays



# What Has Been Achieved?

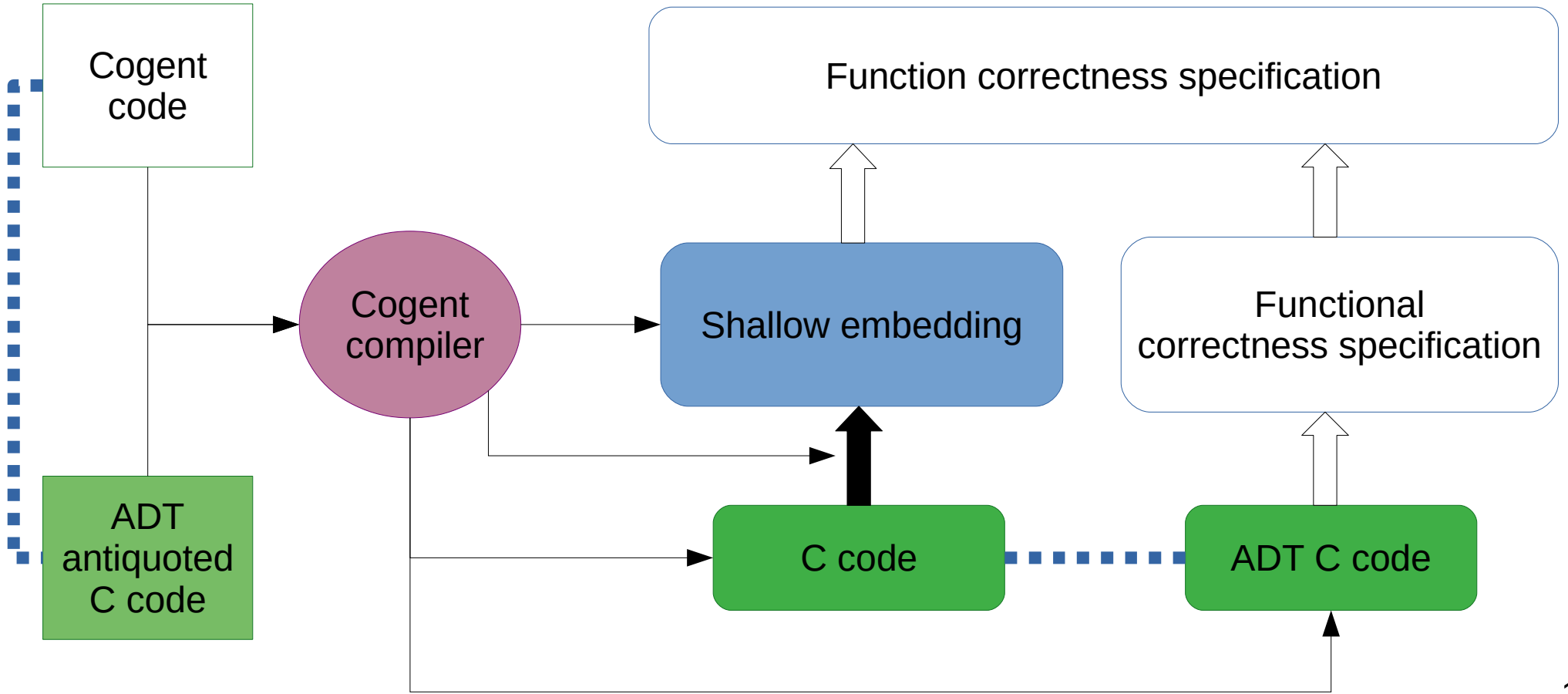
- Defined a functional correctness specification for 32-bit word arrays
- Verified functional correctness for the get, length, put and fold
- Verified frame correctness for get, length, put and fold
- 32-bit word array proofs also work with slight modifications for other standard word lengths

# Methodology

# Functional Correctness Specification

- 2 Choices:
  - 1) Define an abstraction from word arrays to an abstract data type (Isabelle/HOL lists)
  - 2) Define high level HOL properties for word arrays

# Why Option 1?



# Frame Constraints for C Pointers

- Inertia

- $p \notin w_i \cup w_o \Rightarrow \mu_i(p) = \mu_o(p)$

- Leak freedom

- $p \in w_i \wedge p \notin w_o \Rightarrow \mu_o(p) = \perp$

- Fresh allocation

- $p \notin w_i \wedge p \in w_o \Rightarrow \mu_i(p) = \perp$

Pointer :  $p$

Input writable pointer set :  $w_i$

Output writable pointer set :  $w_o$

Input heap function :  $\mu_i$

Output heap function :  $\mu_o$

# Generate the C files

```
$ty:(WordArray a) $id:wordarray_put2($ty:(WordArrayPutP a) args)
{
    if (likely(args.idx < (args.arr->len)) {
        args.arr->values[args.idx] = args.val;
    }

    return args.arr;
}
```

```
WordArray_u32 *wordarray_put2_0(t2 args)
{
    if (__builtin_expect(!(args.idx < args.arr->len), 1))
        args.arr->values[args.idx] = args.val;
    return args.arr;
}
```

# Results

# Bug Discovery

```
struct WordArray_u32 {  
    int len;  
    u32 *values;  
};
```



# Points of Interest

# Pointer Arithmetic

- Accessing the  $i^{\text{th}}$  element in an array
  - $p+k \times i$
- Issues
  - $a \neq b \Rightarrow p+k \times a \neq p+k \times b$  ?

# Word Arithmetic

- $k \times a \equiv k \times b \pmod{n} \Rightarrow a \equiv b \pmod{n} ?$
- Only if  $k$  is coprime with  $n$ 
  - $2 \times 2 \equiv 2 \times 6 \pmod{8}$

# What's the Issue

- 32-bit words are 4 bytes in size
- Word arithmetic in a 32-bit architecture modulo  $2^{32}$
- Overflow can allow 2 distinct indices to have the same addresses in memory
- Assume that 4 times the length of a 32-bit word array is less than  $2^{32}$

# Frame Constraints

- Different memory models
  - Cogent has a single heap
  - Autocorres abstracts the heap as different typed heaps
- Which pointers are writeable?
- Should a Cogent pointer equate to single C pointer or to a set of C pointers?

# Conclusions and Future Work

# Specification

- Functional correctness specification for 32-bit word arrays can be reused for arbitrary word length arrays (with a few tweaks)

# Proof Techniques

- Most of the proofs can be reused for arbitrary word length arrays (with a few tweaks)



# Frame Constraint Satisfiability

- Develop a tactic to automatically extract the writeable pointer sets and to define the frame constraints for each type